Signetics 2650 Microprocessor application memos currently available:

AS50	Serial Input/Output
AS51	Bit and Byte Testing Procedures
AS52	General Delay Routines
AS53	Binary Arithmetic Routines
AS54	Conversion Routines
SP50	2650 Evaluation Printed Circuit Board Level System (PC1001)
SP51	2650 Demo Systems
SP52	Support Software for use with NCSS Timesharing System
SP53	Simulator, Version 1.2
SP54	Support Software for use with the General Electric Mark III Timesharing System
SS50	PIPBUG
SS51	Absolute Object Format (Revision 1)
MP51	2650 Initialization
MP52	Low Cost Clock Generator Circuits
MP53	Address and Data Bus Interfacing Techniques

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INTRODUCTION

Interfacing a microprocessor to peripheral devices is an important part of a total microcomputer system design. The characteristics of the interface depend to a large extent on total system requirements and other factors such as CPU loading and data speed. The use of interrupts and/or DMA structures also have an impact on the system input/output structure. The design of an I/O interface is not limited to hardware, and hardware/software trade-offs must be considered.

This applications memo examines the use of the 2650's versatile set of I/O instructions and the interface between the 2650 and I/O ports. Interrupt and DMA-controlled I/O are not discussed. A number of application examples for both serial and parallel I/O are given. Several types of input, output, and bidirectional interface devices are also examined.

Basic I/O Structure of the 2650

The 2650 is equipped with extensive and versatile input and output facilities. It can perform both single bit input/output and 8-bit parallel input/output.

The single bit input and output, called Sense (pin 1) and Flag (pin 40), are associated with the Program Status Word Upper (PSWU). The Flag output always reflects the value of bit 6 of the PSWU, while bit 7 of the PSWU always reflects the value of the Sense input signal. The Sense and Flag signals can be monitored and controlled with the PSW instructions.

Parallel I/O can be accomplished using the extended or non-extended read and write instructions. The extended and non-extended types are distinguished by the state of the E/$\overline{NE}$ output of the 2650.

The non-extended I/O instructions are single-byte instructions which accomplish a 1-byte data transfer into or out of the 2650. They also control the state of the D/$\overline{C}$ output, which can be used as a 1-bit device address in small systems.

The extended I/O instructions are 2-byte instructions. When executing extended I/O instructions, the second byte of the instruction is output on the lower 8 bits of the address bus (ADR0-ADR7). This information is normally used as an I/O device address to select 1 of up to 256 input or output devices, but may also be used to output control or status signals.
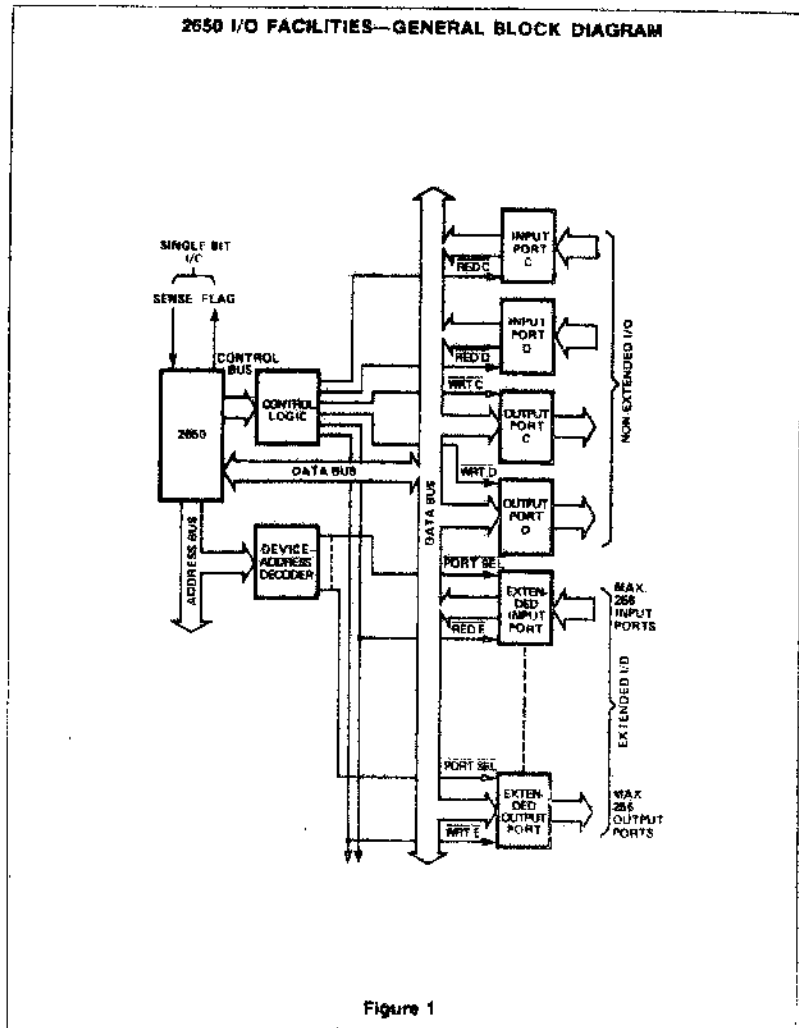

2650 I/O FACILITIES—GENERAL BLOCK DIAGRAM

Figure 1

Parallel I/O operations may use any CPU register as the data source or destination. This offers significant flexibility in writing I/O software, because there is not a single accumulator register to create a "bottleneck" in the data flow. The functional block diagram in Figure 1 illustrates the various I/O facilities.

I/O As Part of the Memory Address Space

The 2650 user may choose to transfer data into or out of the processor using the memory control signals. The advantage of this technique is that the data can be read or written by the program with memory load and store instructions, and data may be directly operated upon with logical and

arithmetic instructions. The memory referencing instructions can take advantage of the flexible addressing modes provided by the 2650, such as indexing and indirect addressing. A possible disadvantage of this method is that it may be necessary to decode more address lines to determine the device address than with the other I/O facilities.

To make use of this technique, the designer must assign memory addresses to I/O devices and design the device interfaces to respond to the same signals as memory.

I/O Interface Signals

Table I summarizes the state of the 2650 I/O interface signals for the various methods of I/O which are available.

SERIAL I/O USING THE SENSE INPUT AND FLAG OUTPUT

One of the I/O capabilities of the 2650 is provided by the sense input and flag output. The sense and flag pins may be used for single-bit input or output of status or control information. They can also be used to implement a serial data communications channel. Two examples of this application are given below.

Asynchronous Serial Communications Port

In applications where a serial type of terminal (like a teletypewriter) must be connected to the microcomputer system, the sense pin and flag pin can be used to interface with the terminal. The basic character format for asynchronous serial I/O is shown in Figure 2.

A number of parameters of this character format and the transmission speed, are different for various types of terminals. The variable parameters are:

Baud rate (bits per second): 110, 150, 300, 600, 1200, 2400, 4600, and 9600 baud.

Number of bits per character: 5, 6, 7, or 8 bits

Parity mode: even, odd, and no parity

Number of stop bits: 1 or 2

The control of the sense and flag pins for asynchronous serial I/O, with the appropriate parameters and baud rate, can be done completely with software. The hardware involved is limited to a simple line driver and receiver circuit which may be either an RS-232 interface or a 20mA current loop interface. The interface hardware is shown in Figure 3.

The software necessary to accomplish the serial I/O for a full-duplex line can be divided into 3 parts:

• The start bit detection and verification. After each start bit detection the start-bit level is verified for a low level at time intervals of 1/6 of 1 bit time. This prevents false start-bit recognition caused by line noise.

• The sampling of the data bits at the mid-bit time, echoing the data bit to the flag output, and loading the data bit into a CPU register.

• The input echo and check of parity bit and stop bits.

A timing diagram showing the start bit sampling and the bit echo appears in Figure 4.

TYPE OF I/O OPERATION	OPREQ	M/IO	R/W	ADR0-ADR7	ADR13 (E/NE)	ADR14 (D/C)
Sense (Input)	X	X	X	X	X	X
Flag (Output)	X	X	X	X	X	X
Extended Read	H	L	L	Second Byte of Instruction	H	X
Extended Write	H	L	H		H	X
Non-Extended Read C	H	L	L	X	L	L
Non-Extended Read D	H	L	L	X	L	H
Non-Extended Write C	H	L	H	X	L	L
Non-Extended Write D	H	L	H	X	L	H
Memory I/O Read	H	H	L	ADR0-ADR7	ADR13	ADR14
Memory I/O Write	H	H	H	ADR0-ADR7	ADR13	ADR14

X = Don't Care

Table 1   I/O INTERFACE SIGNAL STATE

BASIC CHARACTER FORMAT FOR ASYNCHRONOUS I/O

Figure 2

SERIAL I/O HARDWARE INTERFACES
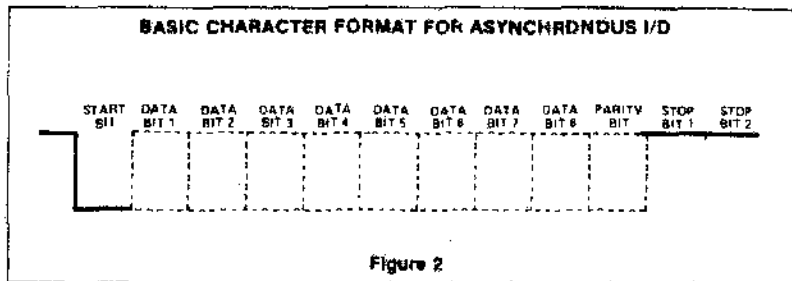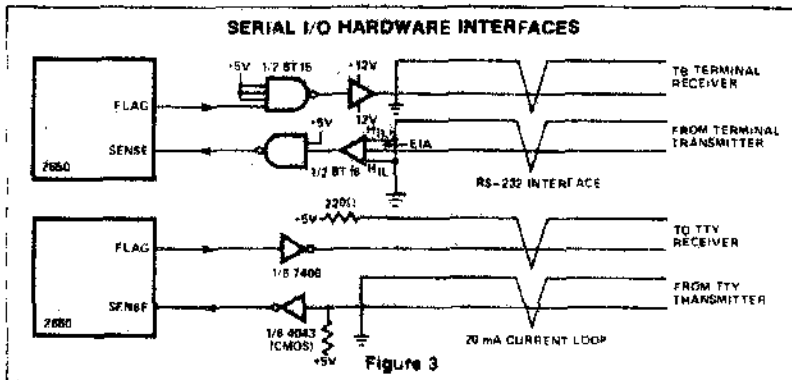
Figure 3
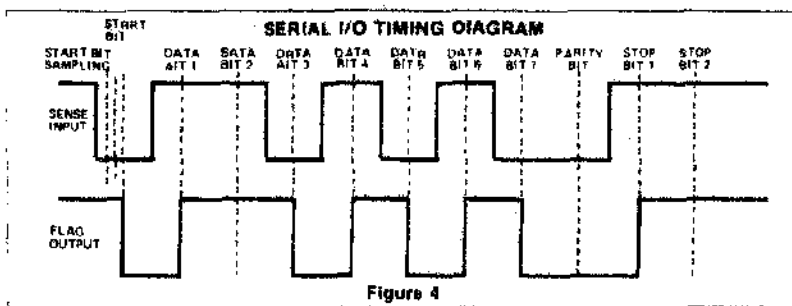

SERIAL I/O TIMING DIAGRAM

Figure 4

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Three examples of the serial I/O routine with different speeds and parameters are presented in Figures 5 through 9. The bit and sample delay numbers (hexadecimal) in the definition listing (Figure 6) are for a CPU clock frequency of 1MHz. The hexadecimal delay numbers for a frequency of 1.25MHz are given in Table II. This table also lists the number of BDRR,R0 instructions that are necessary in the "bit delay and echo subroutine" to count cycles for the appropriate baud rate.

The examples of figures 7, 8, and 9 have the following parameters:

Figure 7: 110 baud, 7 data bits, even parity and 1 stop bit.
Figure 9: 600 baud, 7 data bits, odd parity and 2 stop bits.
Figure 9: 2400 baud, 8 data bits, no parity and 1 stop bit.

The serial I/O routine uses 4 CPU registers (1 bank and R0) and affects 7 of the Program Status Word bits; namely, Sense, Flag, Overflow, Carry, Interdigit Carry, and the 2 Condition Code bits. The program also uses 1 level of the return address stack.

A parity error will set the Overflow bit, and a framing error (wrong stop bit level) will set the Interdigit Carry bit. At the end of the routine, the input character is stored in register R2.

FLOWCHART OF THE SERIAL I/O ROUTINE

Figure 5

SERIAL I/O PARAMETER DEFINITIONS

SERIAL I/O ASSEMBLY LISTING—EXAMPLE 1

Figure 6

Figure 7

SERIAL I/O ASSEMBLY LISTING

EXAMPLE 2

```
TWIN ASSEMBLER VER 1.0                          PAGE 0003

LINE ADDR OBJECT E SOURCE

0073         ****************************************************
0074         * EXAMPLE 2 FULL DUPLEX (BIT BY BIT ECHO) 600 BAUD
0075         * 7 DATA BITS, ODD PARITY AND 2 STOP BITS
0076         *
0077 0600        ORG   H'0600
0078 0600 7640 STRT PPSU  F          SET FLAG TO SWITCH OFF THE LINE
0079 0602 7528      CPSL  CMP+L+IDC
0080 0604 12   TEST SPSU             WAIT FOR START BIT
0081 0605 1970      BCTR,N TEST
0082 0607 0604 LODI,R2 H'03          SET R2 TO NUMBER OF SAMPLES
0083 0609 05DC SAMP LODI,R1 SDE6     SET R1 TO SAMPLE DELAY
0084 060B F97E      BDRR,R1 1
0085 060D 12        SPSU             TEST FOR START BIT VALIDITY
0086 060E 1A74      BCTR,N TEST      IF NOT VALID, GO BACK TO TEST
0087 0610 FA77      BDRR,R2 SAMP
0088 0612 0780      LODI,R3 0P       SET R3 TO ODD PARITY MODE
0089 0614 0607      LODI,R2 0R7      SET R2 TO NUMBER OF DATA BITS
008A 0616 7440      CPSU  F          GENERATE START BIT
0091 0618 51   BITS  RRR,R1
0092 0619 201A      BSTR,UN BDLY     GO TO DELAY AND ECHO ROUTINE
0093 061B FA70      BDRR,R2 BITS     TEST FOR NUMBER OF DATA BITS
0094 061D 01        LODZ  R1
0095 061E C2        STRZ  R2         LOAD R2 WITH CHARACTER
0096 061F 3B14 PRRF BSTR,UN BDLY
0097 0621 9760      BCFR,N STOS
0098 0623 7704      PPSL  OVF        IF WRONG PARITY, SET OVF
0099 0625 9700 STOS LODI,R3 4       CLEAR R3
0100 0627 3000      BSTR,UN BDLY
0101 0629 1A02      BCTR,N SP02      TEST STOP BIT LEVEL
0102 062A 7720      PPSL  IDC        IF WRONG, SET IDC BIT
0103 062C 9700 STO2 LODI,R3 0       CLEAR R3
0104 062E 3004      BSTR,UN BDLY
0105 0631 1A   EKI1 BCTR,N          TEST STOP BIT 2 LEVEL
0106 0632 7720      PPSL  IDC        IF WRONG, SET IDC BIT
0107 0634 17   EKI2 RETC,UN
0108         *
0109         *****************************************************
0110         * BIT DELAY AND ECHO SUBROUTINE
0111         *
0112 0675 0600 BDLY LODI,R0 BRDE     SET R0 TO BIT DELAY NUMBER
0113 0637 FB7E      BDRR,R0 $
0114 0639 12        SPSU             TEST DATA BIT LEVEL
0115 063A 9704      BCFR,N ONE
0116 063C 7440      CPSU  F          IF LOW, ECHO A ZERO
0117 063E 1B04      BCTR,UN BFFL
0118 0640 7640 ONE  PPSU  F          IF HIGH, ECHO A ONE
0119 0642 C540      STRR,R1 $        INSERT DATA BIT INTO R1
0120 0644 23   BFFL LODR  R3
0121 0645 C3        STRZ  R3         DO PARITY CHECK
0122 0646 17        RETC,UN
0123         *
0124 0000      END   $

TOTAL ASSEMBLY ERRORS = 0000
```

EXAMPLE 3

```
TWIN ASSEMBLER VER 1.0                          PAGE 0003

LINE ADDR OBJECT E SOURCE

0073         *****************************************************
0074         * EXAMPLE 3 FULL DUPLEX BIT BY BIT ECHO, 2400 BAUD
0075         * DATA BITS, NO PARITY AND 1 STOP BIT
0076         *
0077 0600        ORG   H'0600
0078 0600 7640 STRT PPSU  F          SET FLAG TO SWITCH OFF THE LINE
0079 0602 7528      CPSL  CMP+L+IDC
0080 0604 12   TEST SPSU             WAIT FOR START BIT
0081 0605 1970      BCTR,N TEST
0082 0607 0604 LODI,R2 H'03          SET R2 TO NUMBER OF SAMPLES
0083 0609 05DC SAMP LODI,R1 SDE6     SET R1 TO SAMPLE DELAY
0084 060B F97E      BDRR,R1 1
0085 060D 12        SPSU             TEST FOR START BIT VALIDITY
0086 060E 1A74      BCTR,N TEST      IF NOT VALID, GO BACK TO TEST
0087 0610 FA77      BDRR,R2 SAMP
0088 0612 0607      LODI,R2 0R7      SET R2 TO NUMBER OF DATA BITS
0089 0614 7440      CPSU  F          GENERATE START BIT
008A 0616 51   BITS  RRR,R1
0091 0617 201A      BSTR,UN BDLY     GO TO DELAY AND ECHO ROUTINE
0092 0619 FA70      BDRR,R2 BITS     TEST FOR NUMBER OF DATA BITS
0093 061D 01        LODZ  R1
0094 061E C2        STRZ  R2
0095 061F 3B14 STOP LODI,R3 0780     CLEAR R3
0096 0621 9760      BCFR,UN BDLY
0097 0623 1A   EKI1 BCTR,N           TEST STOP BIT LEVEL
0098 0624 7720      PPSL  IDC        IF WRONG, SET IDC BIT
0099 0626 17   EKI2 RETC,UN
0100         *
0101         *****************************************************
0102         * BIT DELAY AND ECHO SUBROUTINE
0103         *
0104 0625 0420 BDLY LODI,R0 BRDE     SET R0 TO BIT DELAY NUMBER
0105 0627 FB7E      BDRR,R0 $
0106 0629 12        SPSU             TEST DATA BIT LEVEL
0107 062A 9704      BCFR,N ONE
0108 062C 7440      CPSU  F          IF LOW, ECHO A ZERO
0109 062E 1B04      BCTR,UN BFFL
0110 0630 7640 ONE  PPSU  F          IF HIGH, ECHO A ONE
0111 0632 C540      STRR,R1 BFE      INSERT DATA BIT INTO R1
0112 0634 C3   BFFL STRZ  R3
0113 0635 17        RETC,UN
0114         *
0115 0000      END   $

TOTAL ASSEMBLY ERRORS = 0000
```

Figure 8　　　　　　　　　　　　　　　　**Figure 9**

Data String Output

A typical application for the flag output is a data string output. The advantage of this output method is that it can provide a large number of output bits with little address or control logic decoding. For example, this method can be used to output data for an array of numeric displays, single bit indicators, or column drivers of a parallel numeric printer. An example of the hardware required to implement this type of output channel is given in Figure 10.

Here, the Address 14 output is used as a data strobe signal. However, the data strobe signal could also be built up by decoding more address bits so that the system memory size would not be limited to 16K bytes as in this example.

A listing of the program required is given in Figure 14. The data is assumed to be located in the system's RAM as illustrated in Figure 11.

The least-significant bit of the least-significant byte will be output first. The table length (TLEN) and the number of bits per byte (BPW) can be adapted as necessary by software modifications. The date strobe pulse on output ADR14 is generated by doing the dummy instruction STRA,R0 to address H'4000'.

BAUD RATE	SAMPLE DELAY NUMBER AT 1.25MHz	BIT DELAY NUMBER AT 1.25MHz	NUMBER OF BDRR,R0 INSTRUCTIONS AT 1.25MHz	NUMBER OF BDRR,R0 INSTRUCTIONS AT 1MHz
110	D0	E5	5	4
300	4A	C5	2	2
600	24	DE	1	1
1200	11	6A	1	1
2400	07	30	1	1

Table 2 BIT DELAY PROGRAM CONSTANTS
AT A CLOCK FREQUENCY OF 1.25MHz (HEXADECIMAL)

INTERFACE DIAGRAM FOR DATA STRING OUTPUT

Figure 10

DATA ORGANIZATION FOR DATA STRING OUTPUT

Figure 11

TIMING DIAGRAM OF DATA STRING OUTPUT ROUTINE

TIMES GIVEN FOR A CLOCK FREQUENCY OF 1 MHz

Figure 12

FLOWCHART OF DATA STRING OUTPUT ROUTINE	ASSEMBLY LISTING OF DATA STRING OUTPUT ROUTINE

Figure 13

Figure 14

PARALLEL INPUT/OUTPUT

The 2650 instruction set contains the following six input/output instructions

		NO. BYTES
WRTC, RX	Write Control	1
REDC, RX	Read Control	1
WRTD, RX	Write Data	1
REDD, RX	Read Data	1
WRTE, RX DEVA	Write Extended	2
REDE, RX DEVA	Read Extended	2

The control signals generated by each I/O instruction simplify the interface circuitry required to generate I/O selection and timing signals. A low-cost control signal interface with related timing is shown in Figures 15 and 16.

When using standard TTL and 8T series I/O ports, the I/O operations can be done without slowing down the system. In this case the OPACK input could be controlled directly for all I/O operations.

Non-Extended I/O

The single-byte I/O instructions of the 2650 are referred to as non-extended I/O. In small systems with only two 8-bit input ports and two 8-bit output ports, this I/O facility requires a minimum of hardware interfacing between the CPU and I/O ports. The signals WRTC, WRTD, REDC, and REDD generated by the control logic decoder in Figure 15 can be used directly as output port clock pulses and input port enable signals, respectively.

Sequential I/O With Non-Extended I/O Instructions

In systems where a larger number of devices must be serviced in sequence, the use of a simple 8-bit output port can offer considerable savings in software. Normally the devices could be serviced with extended I/O instructions. However, since the device address is the second byte in this type of instruction, a series of data fetch and I/O instructions would be required to service the devices in sequence.

With an 8-bit output port functioning as a device address register, the device address can be modified under software control. In this way a simple program loop can serve up to 8 I/O ports by rotating a single '1' through a CPU register that is output as a device address. This I/O addressing technique may also be used advantageously in systems where I/O operation requests are detected by software polling. A functional block diagram of this technique is shown in Figure 17.

Extended I/O

There are 2 extended I/O instructions in the 2650 instruction set. In these 2-byte instructions the first byte specifies the operation code and the data source or destination register in the CPU. The second byte provides an 8-bit device address code that is output on the 8 least-significant bits of the address bus, ADR0 through ADR7.

The control signal decoding diagram (Figure 15) can be simplified for systems using only extended I/O, as shown in Figure 18. The timing diagram of Figure 16 also applies to this decoding technique.

Device Address Decoding Schemes

For extended I/O it is necessary to decode the address lines ADR0 through ADR7 in order to generate appropriate port selection signals. The choice of an address decoding scheme depends on factors such as total

CONTROL SIGNAL INTERFACE USING THE 74(L) S138 DECODER

Figure 15

I/O INTERFACE TIMING DIAGRAMS (Figure 15)

Figure 16

I/O requirements, the type of I/O ports used, and the total system configuration.

In principle, there are 2 basic methods of device address decoding. One method is the use of hardwired logic in which the device address is fixed; the other is a hardware programmable method in which the device addresses are individually set with jumpers or switches. Some examples of these methods are given in Figures 19 and 20.

In many applications a combination of these 2 methods is used. In addition the control logic can be implemented as an integral part of the device address decoding. An example is shown in Figure 21.

Memory Mapped I/O

In memory mapped I/O, the I/O devices are treated as memory locations. An advantage of this technique is that all memory referencing instruction types (store, load, arithmetic, logical, etc.) can be used directly for I/O data. Device address decoding is not necessarily more complex than for normal extended I/O, since all I/O addresses could be located in a specific address block. Of course, this technique can only be used in systems which do not use the full memory address space for programs. A diagram of the I/O control logic, using the ADR14 output to discriminate between memory and I/O operations, is given in Figure 22. The device address decoding methods described earlier can also be applied to memory mapped I/O.

SEQUENTIAL I/O TECHNIQUE

Figure 17

SIMPLIFIED CONTROL LOGIC WHEN USING EXTENDED I/O ONLY

(A) Using 1-of-8 Decoder

(B) Using Logic Gates

Figure 18

SOME POSSIBLE TECHNIQUES FOR DEVICE ADDRESS DECODING

(A) Each address line selects one device (maximum of eight)

(B) Using a 1-out-of-16 decoder

Figure 19

HARDWARE PROGRAMMABLE DEVICE ADDRESS DECODERS

(A) Using Exclusive-OR Gates

(B) Using Comparators

Figure 20

COMBINED CONTROL LOGIC AND ADDRESS DECODING

INPULSE AND OUTPULSE CAN BE USED DIRECTLY TO ENABLE INPUT PORTS AND TO CLOCK OUTPUT PORTS

5 JUMPERS FOR PROGRAMMING

Figure 21

SINGLE POINT CONTROL

In many applications, the capability to set, clear, or test a single output point selected from a large number of output points is required. Designs of this type can be implemented using the 2650 I/O instructions. When used as described below, the WRTE, WRTC, and WRTD instructions become "set/clear single-bit" instructions, while the REDE instruction becomes a "test single-bit" instruction.

Single Bit Output—Direct Address

The write extended instruction can be used to select and set or clear a single output bit. The 2 bytes of the instruction can be interpreted as follows:

BYTE 0

1	1	0	1	0	1	X	X

BYTE 1

S/C	A_6	A_5	A_4	A_3	A_2	A_1	A_0

A_0 through A_6 of the second byte specify the output selected. The S/C bit specifies whether the bit is set or cleared. A typical hardware configuration controlling 64 points is shown in Figure 23. Here, the control line decoding and partial address decoding is done by the 74LS138, which selects one of the eight 9334s. One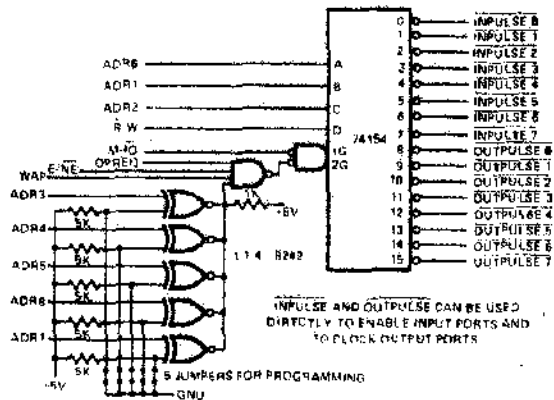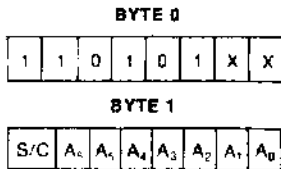 of the 8 latches in the selected 9334 is enabled by ADR0, ADR1, and ADR2 and is either cleared or set, as determined by the value of ADR7.

The XX field in the first byte selects 1 of the 4 available registers and outputs its contents on the data bus. Since this information is not used in this application, the value of XX is not important. However, it could be used to output an 8-bit control or status word in conjunction with the set/clear operation.

Single Bit Output—Indirect Address

If the address of the output to be set or cleared must be determined at program run time, the WRTD and WRTC instructions can be used. The address of the output bit is first loaded into one of the 2650 registers. A WRTD, Rx instruction is then issued if the bit is to be set, and a WRTC, Rx instruction is issued if the bit is to be cleared. The bit select is output on the data bus, and the D/C̄ output carries the set/clear information. The hardware implementation can be the same as shown in Figure 23 except that ADR0-ADR5 are replaced by DBUS0-DBUS5, and ADR7 is replaced by D/C̄.

Figure 22

I/O CONTROL SIGNAL GENERATION FOR MEMORY MAPPED I/O

Figure 23

SIXTY-FOUR SINGLE BIT OUTPUTS USING THE 9334

Single Bit Input

One way of doing single bit input uses the techniques described earlier. The address of the bit that is to be tested is loaded into one of the 2650 registers and output to an 8-bit latch using an extended or non-extended write instruction. The latch output is decoded to select the desired bit, which is then applied to the Sense input pin. The 2650 Program Status Word instructions can then be used to test the state of the Sense input and to take appropriate program action.

The technique described above must be used if "indirect" bit addressing is required. If this is not a requirement, a more efficient implementation can be accomplished using the extended read instruction. This technique makes use of the fact that the 2650 automatically tests the contents of a register every time it is used as the destination of an operation. Thus, when the read extended operation reads data from an input port, the condition code bits in the program status word are set to reflect whether the new

register contents is positive, negative, or zero.

For the single bit input application, the second byte of the RETE, Rx instruction contains the address of the input bit to be tested. This data is applied to a bank of data selectors to select the addressed bit, which is then applied to the most-significant bit of the data bus, DBUS7. Since this is interpreted as the sign bit, the condition code bits in PSL will be set to reflect whether the bit being tested is a one or a zero. A conditional branch instruction can then be used to affect the desired program action. A hardware implementation for 64 inputs is shown in Figure 24. Note that an address latch is not required for this method.

INPUT PORT DEVICES

Gated Input Ports

The simplest form of an input port is the tri-state gate. Figure 25 illustrates the use of the 8T97 high-speed hex tri-state buffer for gated input ports. The 8T97 is non-inverting, and the tri-state control signals enable the buffers in groups of 4 and groups of 2, so that 8-bit ports can be implemented efficiently.

An effective circuit for systems using 8-gated input ports is the 74251 8-to-1 multiplexer, which has tri-state outputs that can interface directly with the data bus. The advantage of this circuit is that no external address decoding logic is needed. A configuration using gated input ports with the 74251 multiplexer is illustrated in Figure 26.

In addition to these 2 configurations, many other input port configurations are possible using standard TTL or Signetics 8T series logic circuits.

Latching Input Ports

Latching input ports may be required to store data from an external device, which is available only momentarily, before the actual input operation to the microprocessor takes place. This type of input port can be realized by connecting TTL-latch or D-type flip-flop circuits, such as the 7475, 74100, or 74175, to the inputs of a gated input port. As illustrated in Figure 27, by using the Signetics 8T10 Quad D-type flip-flop with tri-state outputs, an 8-bit latching input port can be implemented with only 2 packages. The 8T10 is functionally identical to the 74173.
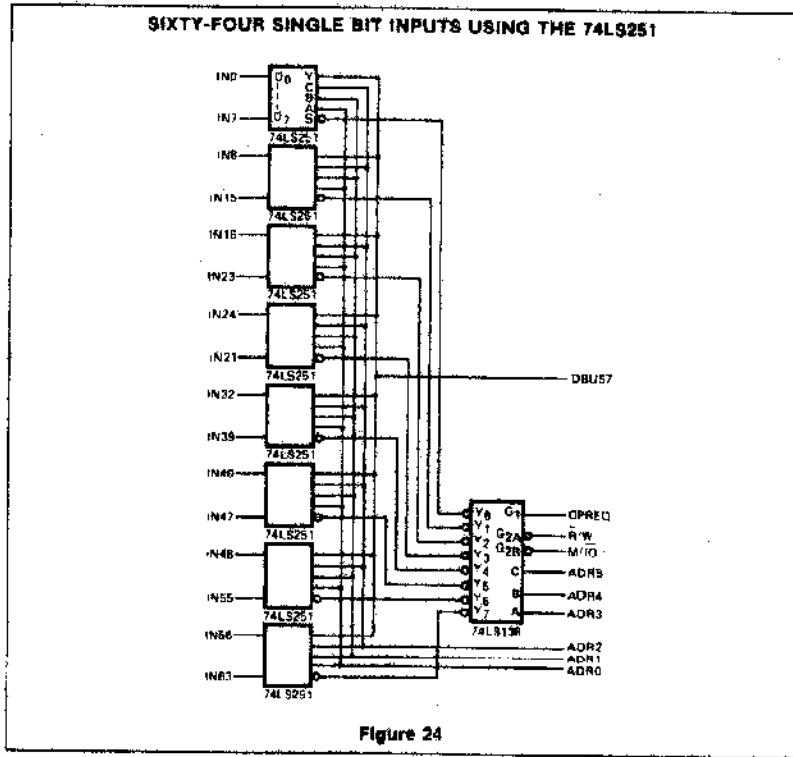
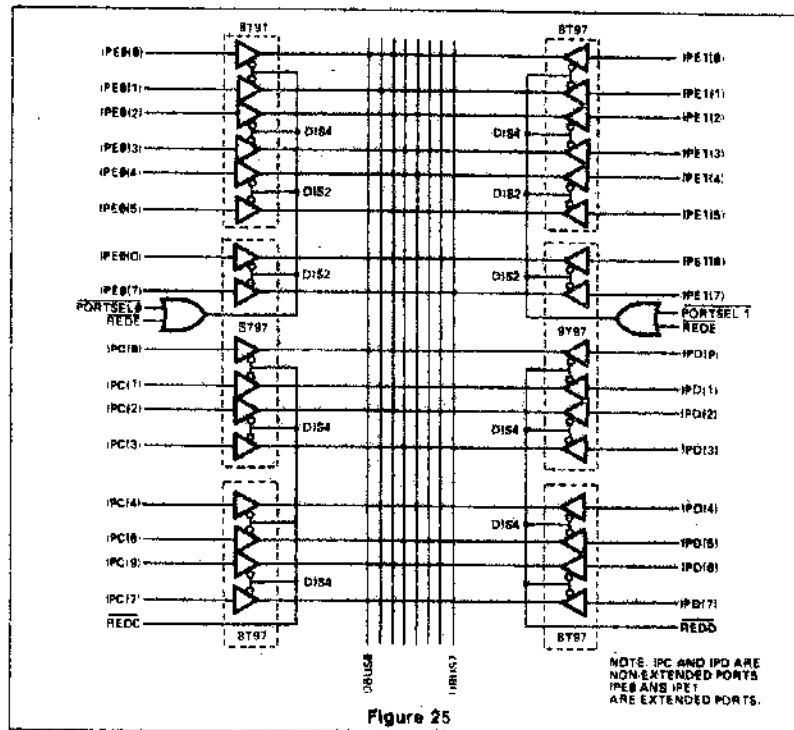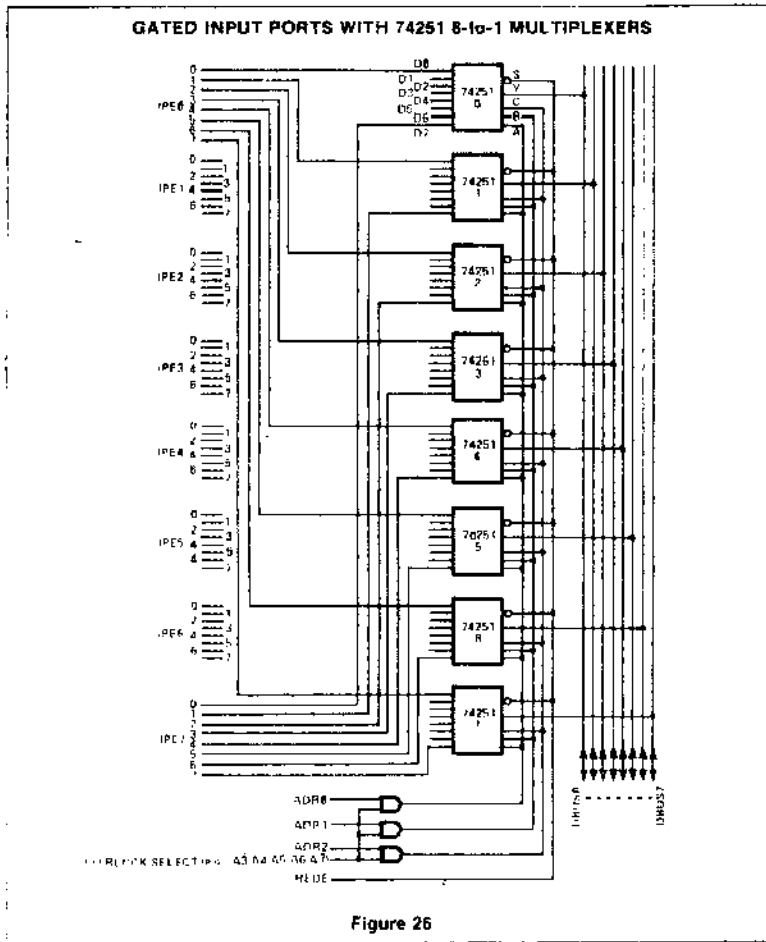**SIXTY-FOUR SINGLE BIT INPUTS USING THE 74LS251**

Figure 24

NOTE: IPC AND IPD ARE NON-EXTENDED PORTS IPE8 AND IPE1 ARE EXTENDED PORTS.

Figure 25

GATED INPUT PORTS WITH 74251 8-to-1 MULTIPLEXERS

Figure 26

LATCHING INPUT PORTS USING THE 8T10 (74173)

Figure 27

OUTPUT PORTS BUILT WITH STANDARD TTL AND 8T SERIES

Figure 28

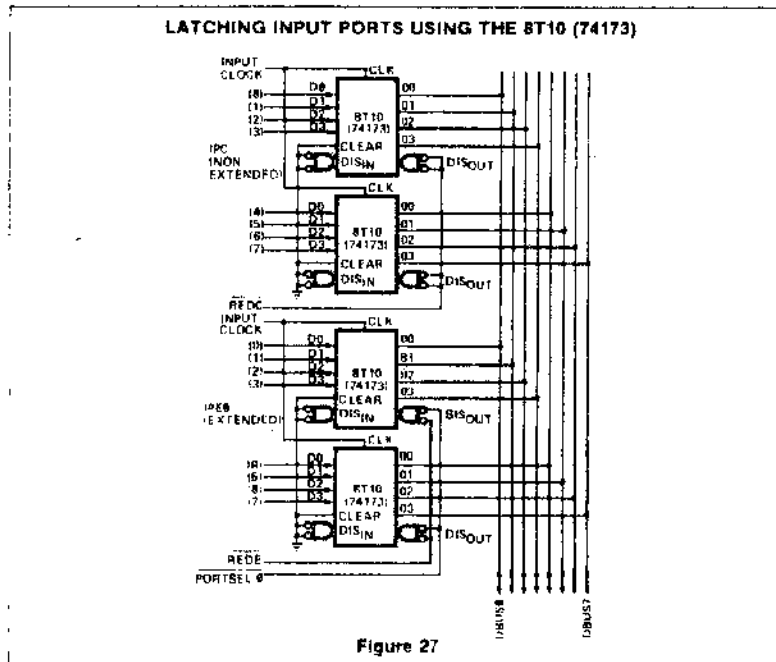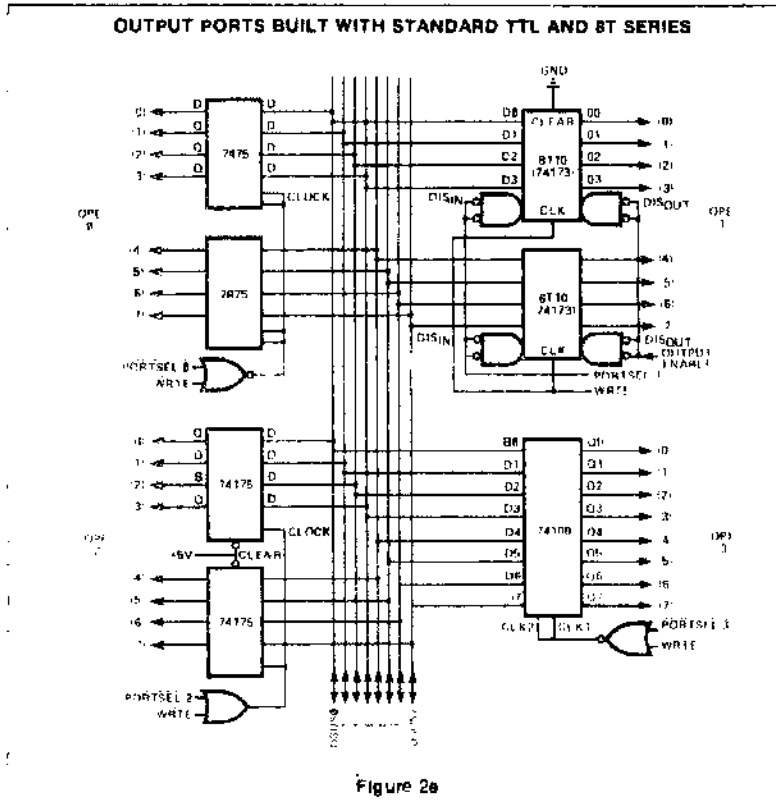## OUTPUT PORT DEVICES

Output ports can be configured with a variety of standard TTL and 8T series flip-flops and registers. Typical circuits include:

9334	Addressable 8-bit latch
7475	Quadruple latch
74100	8-bit latch
74175	Quadruple D-type flip-flop
8T10	Quadruple Q-type flip-flop with tri-state outputs

The 7475 and 74175 both have true and complement outputs. One special feature of the 8T10 is that the outputs may be disabled (placed in a high-impedance output mode) by the device that is connected to this output port. A logic diagram using these circuits for output ports appears in Figure 28.

The 9334 is useful in systems requiring a large number of latched outputs, since a portion of the decoding can be done using the on-chip 3-input decoder. A typical application of this was shown in Figure 23. It is also an efficient circuit for implementing eight 8-bit output ports.

I/O CONFIGURATIONS USING THE 8T31 BIDIRECTIONAL PORT

Functional Description

The 8T31 is an 8-bit bidirectional I/O port consisting of 8 clocked latches with 2 bidirectional I/O buses, each of which has its own control logic. Each bus (A and B) has a read and a write control input, and there is a master enable input for bus B only. The outputs of the latches follow the inputs when the clock is high, and latching will occur when the clock returns low.

The 8T31 is also equipped with a "power-on clear" circuit. If the clock input is held low until the power supply reaches 3.5 volts, the latches will be cleared. There is a logic inversion between bus A and bus B. As a result when the 8T31 is cleared, bus A will have all logic 1 outputs and bus B all logic "0" outputs.

The control functions of the 8T31 are listed in Table III. A functional block diagram and a symbolic diagram of the 8T31 are illustrated in Figures 29 and 30 respectively.
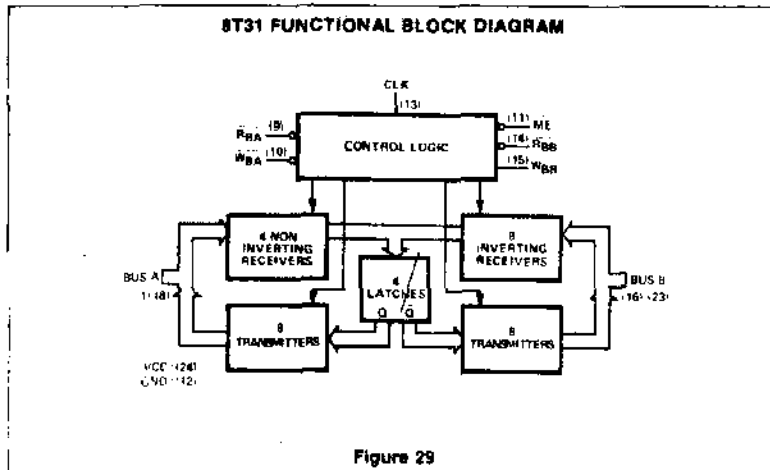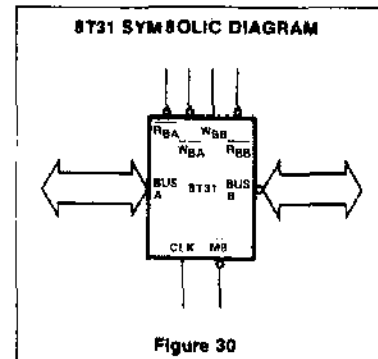
8T31 FUNCTIONAL BLOCK DIAGRAM

Figure 29

8T31 SYMBOLIC DIAGRAM

Figure 30

BUS A			
$\overline{R}_{BA}$	$\overline{W}_{BA}$	CLK	BUS A
X	0	1	WRITE (A→latch)
0	1	X	READ (latch→A)
1	1	X	HI-Z (Tri-state)

BUS B					
$\overline{R}_{BB}$	W_{BB}	$\overline{W}_{BA}$	CLK	$\overline{ME}$	BUS B
X	X	X	X	1	HI-Z
1	0	X	X	0	HI-Z
X	1	0	X	0	HI-Z
0	0	X	X	0	READ (latch→B)
X	1	1	1	0	WRITE (B→latch)

Table 3   8T31 CONTROL FUNCTIONS

As shown in Table III, each bus can operate independently except for the case of writing from both bus A and B. In this case writing from bus A will override any attempt to write from bus B.

8T31 Applications

The control functions of the 8T31 allow it to be used in various microcomputer input output applications. In the I/O system diagram of Figure 31, the 8T31 is used to implement gated input ports, latching input ports output ports, and a bidirectional data bus driver. All I/O ports can be controlled directly with the device select and REDE and WRTE lines coming from device decoders and I/O control logic.

In applications where interfacing is necessary with peripheral devices that need data transfers in two directions, like digital cassettes and data link communication circuits, the 8T31 can be used as a bidirectional I/O port. In this application, the I/O opera-

tion should be requested by interrupt or polling to prevent simultaneous write operations from peripheral and CPU. The bidirectional I/O port concept is illustrated in Figure 32

Implementing an Eight-Bit Flag Register with the 8T31

In many industrial applications, such as process control, single bit inputs and outputs are used to monitor switches and detectors or to drive relays and lamps. A possible solution for such a flag register would be an eight-bit output port and a memory byte reserved as a flag register in the system's RAM. The setting, resetting, or testing of individual bits with this method of implementing a flag register requires many bytes of program memory. The output port and the memory location reserved as a flag register image must be updated after each bit operation

The 8T31 can be used to implement a flag register without the use of a memory byte in the system's RAM. No additional hardware is required and the saving in program memory bytes for flag operations is considerable. A logic diagram of this application is given in Figure 33. Listings of basic software to set, reset, and test individual flags for both positive and negative true outputs are given in Figures 34 and 35

THE 8T31 USED AS A GATED INPUT PORT,
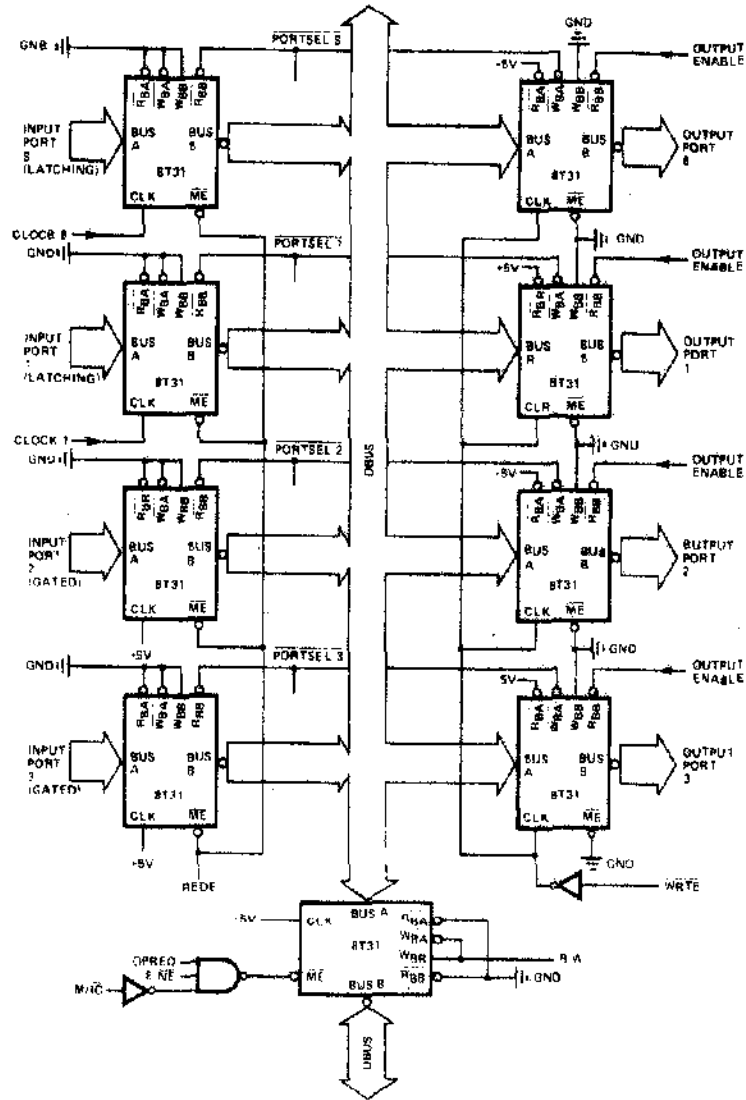LATCHING INPUT PORT, OUTPUT PORT, AND DATA BUS DRIVE

Figure 31